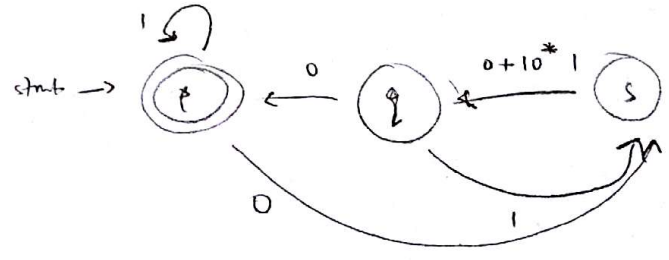


elimination of state r

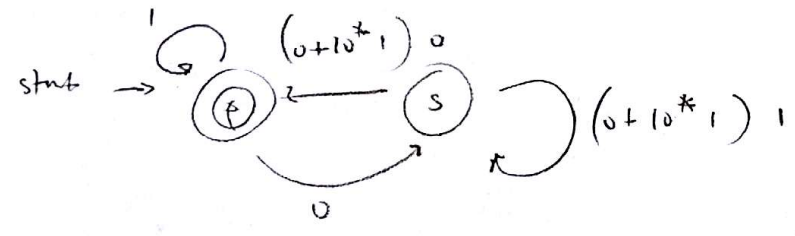
$$(s, q) = 0 + 10^*1$$



elimination of q

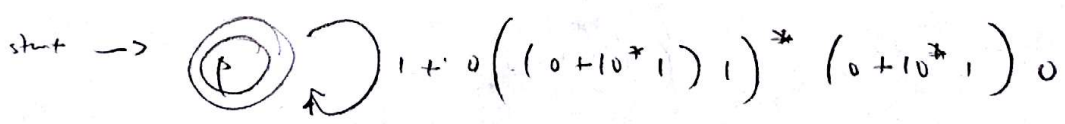
$$(s, p) = p + (0 + 10^*1) \emptyset^* 0 = (0 + 10^*1) 0$$

$$(s, s) = \emptyset + (0 + 10^*1) \emptyset^* 1 = (0 + 10^*1) 1$$



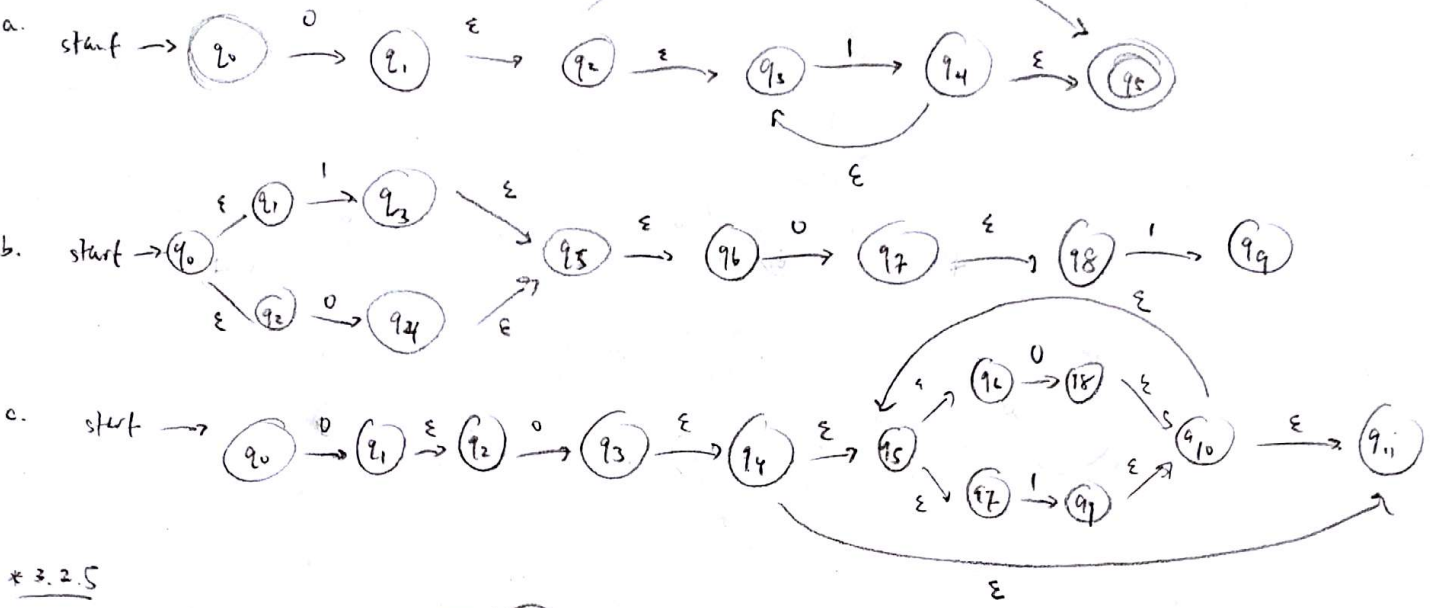
elimination of s

$$(p, p) = 1 + 0 \left((0 + 10^*1) 1 \right)^* (0 + 10^*1) 0$$

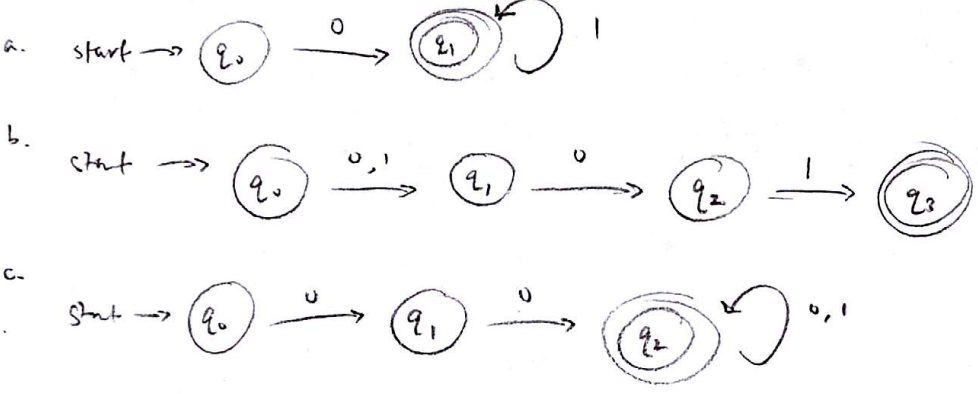


$$= \left(1 + 0 \left((0 + 10^*1) 1 \right)^* (0 + 10^*1) 0 \right)^*$$

*3.2.4



*3.2.5



*3.4.1

d. $R(s+T) = RS + RT$

Suppose a string w is in $R(s+T)$, then $w = xy$ where $x \in R$ and $y \in$ either S or T . If $y \in S$, then $xy \in RS$ and therefore $\in RS + RT$. Conversely, if $y \in T$, then $xy \in RT$ and therefore in $RS + RT$.

Suppose w is in $RS + RT$, then w is either RS or RT . Then for $w = xy$, if $x \in R$ and $y \in S$, then since $y \in S$, it is in $(s+T)$ and conversely for $y \in T$.

e. $(R+S)T = RT + ST$

Suppose a string w is in $(R+S)T$, then $w = xy$ where $x \in$ either R or S and $y \in T$. If $x \in R$, then $xy \in RT$ and therefore $\in RT + ST$. Conversely if $x \in S$, then $xy \in ST$ and therefore $\in RT + ST$.

Suppose $w \in RT + ST$, then $w \in$ either RT or ST . Then for $w = xy$, if $x \in R$ and $y \in T$ then since $x \in R$ then it is in $(R+S)$ and conversely for $x \in S$.

$$g. (\epsilon + R)^* = R^*$$

by substitution of symbols: $(\epsilon + a)^* = a^*$

$$\text{since } \epsilon^* = \epsilon \text{ and } a^* = a^0 + a^1 + a^2 + \dots \\ = \epsilon + a + aa$$

ϵ is already contained by a^*

$$h. (R^* S^*)^* = (R+S)^*$$

by substitution of symbols:

$$(a^* b^*)^* = (a+b)^*$$

It is clear that both sides intuitively represent all strings of a and b .

thus, the two concrete expressions denote the same language so $(R^* S^*)^* = (R+S)^*$ holds.

* 3.4.2

$$b. (RS + R)^* R = R(SR + R)^*$$

True

concrete expression:

$$(ab + a)^* a = a(ba + a)^*$$

by visual inspection, it is clear that both these languages represent strings that start with a and contain any number of a 's or single occurrences of b that end in a .

$$d. (R+S)^* S = (R^* S)^*$$

concrete expression:

$$(a+b)^* b = (a^* b)^*$$

False, proof by contradiction

The string $bb \in L((a+ab)^* b)$ but not $L((a^* b)^*)$. Thus, the

concrete expression does not hold and neither does the statement in general.